

# Interactive constraints computer-aided composition

## ICMC 2017

**Pierre Talbot** Carlos Agon Philippe Esling  
(talbot@ircam.fr)

Institute for Research and Coordination in Acoustics/Music (IRCAM)  
Université Pierre et Marie Curie (UPMC)

19th October 2017

# Computer-aided composition

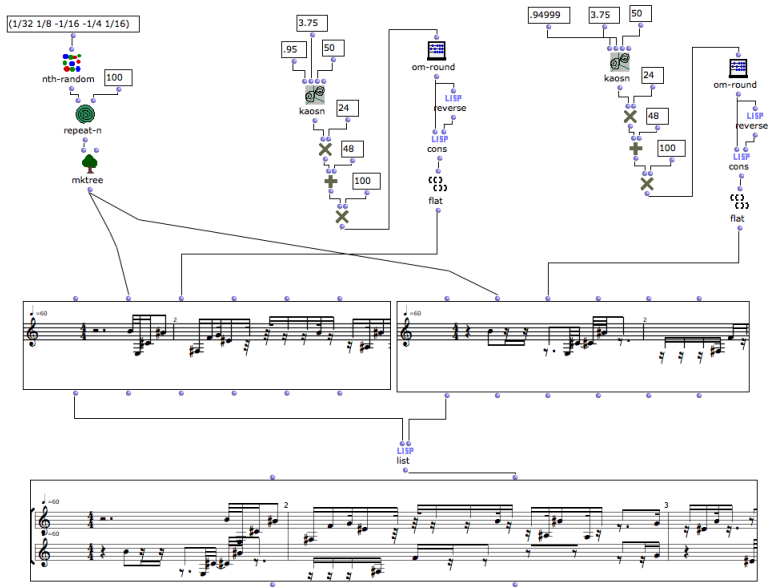
## Goals

- ▶ Delegating tedious computations to the machine.
- ▶ Parametrizing the patch with values to quickly try-and-test.
- ▶ ...

How does the composer interact with the machine?

- ▶ Mostly visual and dataflow programming languages: OpenMusic, PureData, Max,...
- ▶ Functional programming languages for the specifics: Lisp mostly.

# Dataflow: a patch in OpenMusic



# Constraints in computer-aided composition

## Constraint programming

- ▶ Declarative paradigm for solving combinatorial problems.
- ▶ We state the problem and let the system solve it for us.
- ▶ Example: pitches must form a decreasing sequence (from highest to lowest).

Some examples of attempts to add constraints into CAC softwares:

- ▶ PWConstraints on top of PatchWork: constraints over the pitches, grouping the pitches together (modelling aspects).
- ▶ OMCloud on top of OpenMusic is based on a different constraint solving paradigm—local search—aiming at the ease of use.

# Problem

- ▶ CAC softwares extended with constraints work in **black box**: one solution gets out of the box.
- ▶ But constraints are relations, not functions.
- ▶ Therefore, a constraint problem can have zero, one or many solutions.

By functionalizing the constraint process, we miss a key point:

Constraints are useful to describe a class of solutions

but how to work with many solutions?

## Proposal: Interactivity

Experiment with an interactive constraint score editor.

- ▶ Bring the composer at the level of the solving process.
- ▶ He can consciously choose a solution.
- ▶ Development of an interactive search strategy to navigate in the **solution space**.

# Proposal: Spacetime programming

Interactivity and search strategies is a deeper problem: constraint solvers also work in a “black box” mode.

We propose the process calculi **spacetime programming**.  
SP = constraint programming + synchronous paradigm.

## Spacetime programming

- ▶ Synchronous programming for *interactive computing*.
- ▶ A search strategy is viewed as a process: abstraction over the constraint solver.

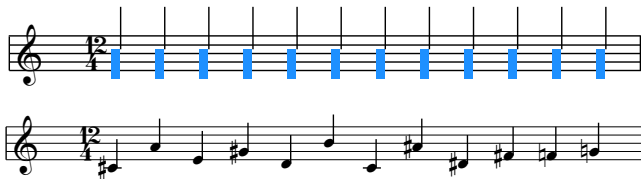
# Menu

- ▶ Introduction
- ▶ Interactivity in solvers
- ▶ Interactivity in CAC
- ▶ Conclusion



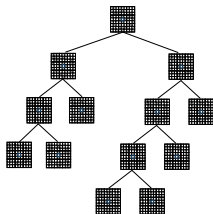
# All-interval series: a MiniZinc model

```
int: n = 12;  
array[1..n] of var 1..n: pitches;  
array[1..n-1] of var 1..n-1: intervals;  
constraint forall(i in 1..n-1)  
    ( intervals [i] = abs(pitches[i+1] - pitches[i]));  
constraint alldifferent ( pitches );  
constraint alldifferent ( intervals );  
  
solve satisfy;
```



# All-interval series: a MiniZinc model

```
int: n = 12;  
array[1..n] of var 1..n: pitches;  
array[1..n-1] of var 1..n-1: intervals;  
constraint forall(i in 1..n-1)  
    ( intervals [i] = abs(pitches[i+1] - pitches[i]));  
constraint alldifferent ( pitches );  
constraint alldifferent ( intervals );  
  
solve satisfy;
```



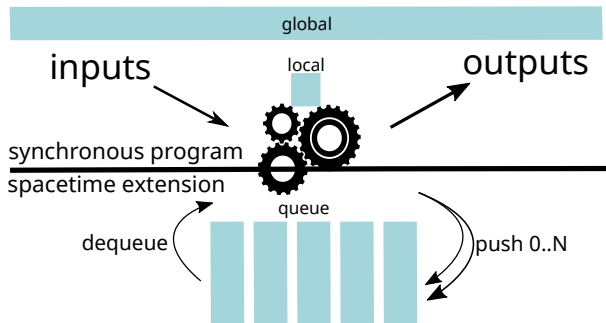
# Synchronous paradigm

- ▶ Invented in the 80s to deal with reactive system subject to many (simultaneous) inputs.
- ▶ Continuous interaction with the environment.
- ▶ Mainly used in embedded systems.



# Spacetime execution scheme

- ▶ The search tree is represented as a queue of nodes.
- ▶ We feed the program with **one node of the tree per instant**.
- ▶ The synchronous program fuels the queue with new nodes.



# Spacetime programming

## Syntax

$\langle p, q, \dots \rangle ::=$	
<i>spacetime</i> <i>Type</i> $x = e$	<b>communication fragment</b> (variable declaration)
<b>when</b> <i>cond</i> <b>then</b> $p$ <b>end</b>	(ask)
$x \leftarrow e$	(tell)
$x.m(\dots)$	(method call)
	<b>synchronous fragment</b>
<b>par</b> $p \parallel q$ <b>end</b>	(parallel composition)
$p ; q$	(sequential composition)
<b>suspend when</b> <i>cond</i> <b>in</b> $p$ <b>end</b>	(suspension)
<b>loop</b> $p$ <b>end</b>	(infinite loop)
<b>pause</b>	(delay)
	<b>search tree fragment</b>
<b>space</b> $p$ <b>end</b>	(branch creation)
<b>prune</b>	(branch pruning)

# Spacetime attribute

## Problem

How to differentiate between variables in internal/global state and those onto the queue?

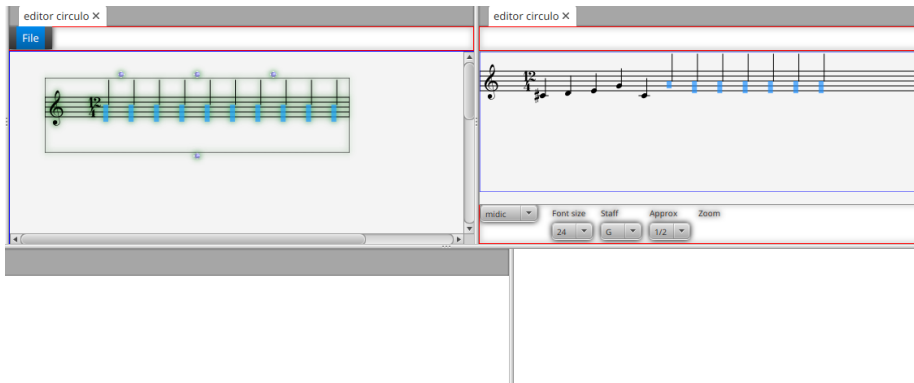
We use a spacetime attribute to situate a variable in space and time.

- ▶ `single_space`: variable **global** to the search tree.
- ▶ `single_time`: variable **local** to one instant.
- ▶ `world_line`: **backtrackable** variable in the queue of nodes.

# Menu

- ▶ Introduction
- ▶ Interactivity in solvers
- ▶ **Interactivity in CAC**
- ▶ Conclusion

# Score editor: overview



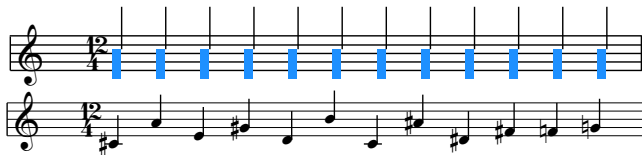
Constraint solving zone for the interactions with the composer.



# A first interactive strategy

The strategy usually implemented in CAC with constraints: stop at each solution. In practice: click on “space” to jump to the next solution.

```
class EachSolution {  
  world_line VStore domains = bot;  
  world_line CStore constraints = bot;  
  proc stop_at_solution =  
    loop  
      par  
        || when domains |= constraints then stop end  
        || pause  
      end  
    end  
  end  
}
```



# Interaction with the composer

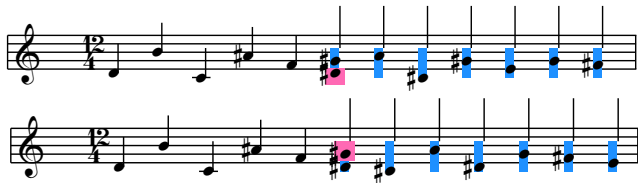
The composer interacts with the search in-between instants.  
The spacetime attributes enable interactions with the search in two main ways: globally or only for the current search path.

```
class PSolver {  
    world_line CStore constraints = bot;  
    single_space CStore cpersistent = bot;  
    ...  
}
```



## Lazily navigating the solution space

The next two scores represent a choice between  $\sharp D$  and  $\sharp G$  on the sixth note:



```
SubSolver<RBinary, Model> left = new SubSolver();
SubSolver<Binary, Model> right = new SubSolver();
single_time L<Boolean> choice = bot;
choice <- top;
par
  || suspend when choice |= true then right.search() end
  || suspend when choice |= false then left.search() end
end
```

# Menu

- ▶ Introduction
- ▶ Interactivity in solvers
- ▶ Interactivity in CAC
- ▶ Conclusion

# Constraints in music

From a computer scientist perspective

- ▶ Probably not for generating music: machine learning methods do it better.
- ▶ Reasoning on a class of scores satisfying some properties.  
Example: we are not forced to write a particular pitch but a class of pitches satisfying some rules.
- ▶ Constraints do not force the composer to make any choice!

# Conclusion

- ▶ Constraints are relational: interactive search helps to use them in this way.
- ▶ To program interactive search strategies, we use spacetime programming.

## Future work

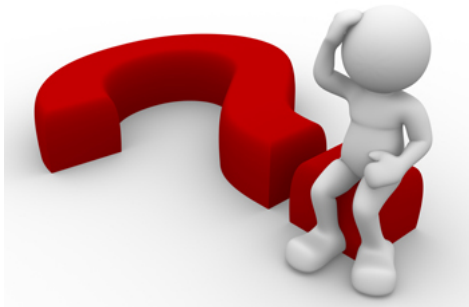
- ▶ Current prototype with AIS only; enabling any MiniZinc model.
- ▶ This would allow composers to try the system and to develop more strategies.

Stay tuned!



[github.com/ptal/bonsai](https://github.com/ptal/bonsai)  
[github.com/ptal/repmus](https://github.com/ptal/repmus)

Thank you for your attention.



Stay tuned!



[github.com/ptal/bonsai](https://github.com/ptal/bonsai)  
[github.com/ptal/repmus](https://github.com/ptal/repmus)